A Semantic-based Middleware for Multimedia Collaborative Applications

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Outline

Introduction

🟓 Middleware

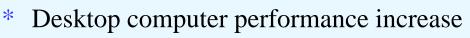
- Objectives
- ➡ Extension of operating systems network services
- Stream synchronization
- Floor control framework
- Protocol for dynamic image transmission
- Experimental results
 - **Conclusions**



Introduction

more

• Large-scale Multimedia Applications



* Internet growth in bandwidth and # of hosts

• A challenging class of applications

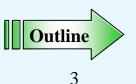
- * Processing power & bandwidth
- * Scalability Imore
- * Heterogeneity (Ethernet/modem, WinNT/Solaris, MPEG/H263)
- * Timely data delivery

• Traditional services

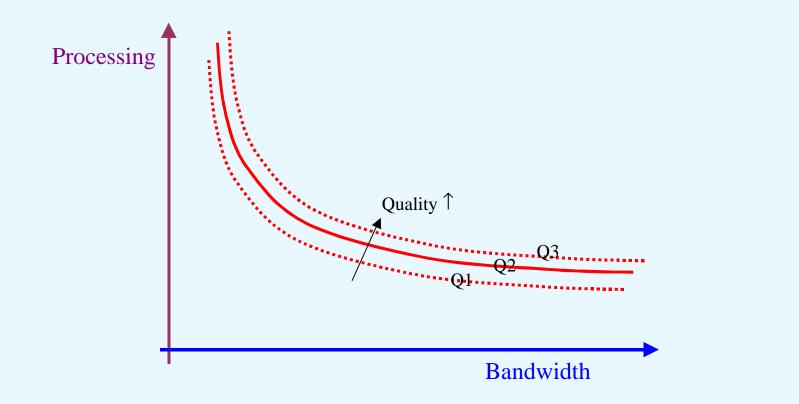
- * Network layer: UDP & TCP (real time was not a concern)
- * Operating systems: Abstractions are not adequate for multimedia.
 - » Example: Real time is not well supported.
- Gap between multimedia requirements and system

services



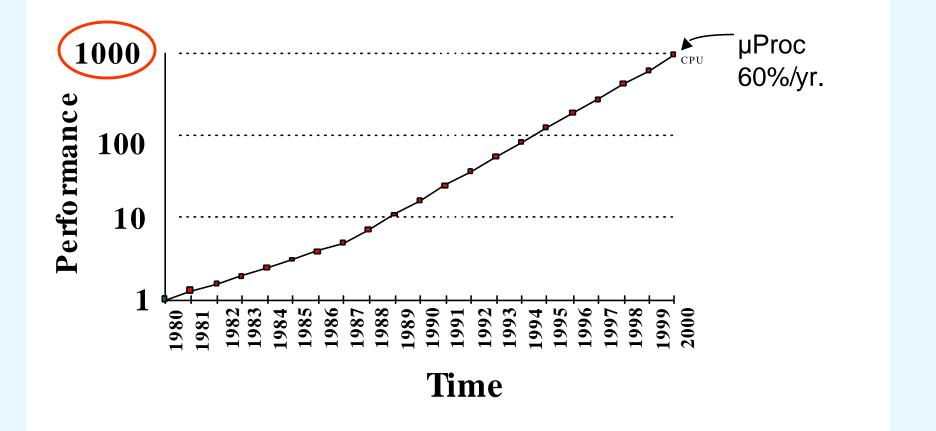


Multimedia Resource Requirements





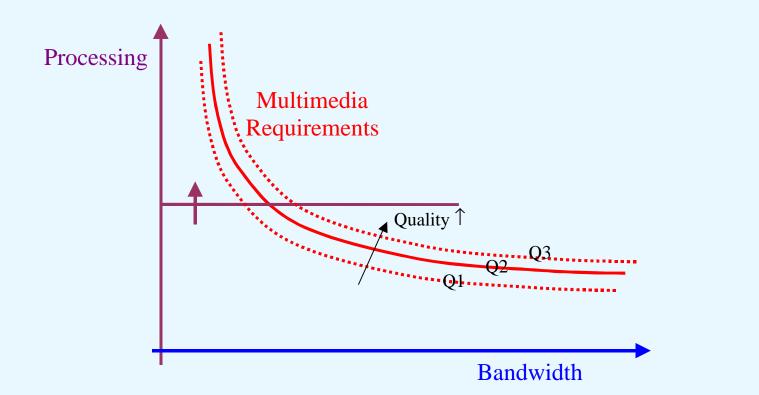
Processor Performance Increase

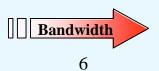


Source: Dr. David Patterson University of Virginia Distinguished Lecture Series, May 19,1998. http://www.cs.berkeley.edu/~pattrsn/talks/Stanford.pdf

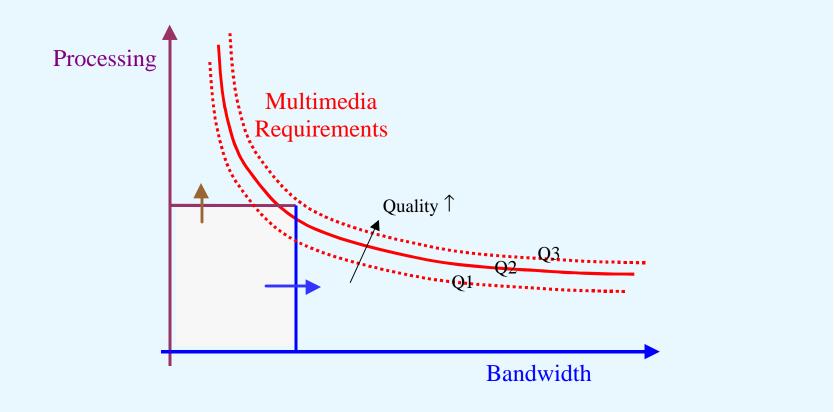
Effect on MM

Multimedia Resource Requirements

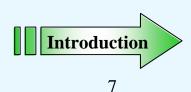




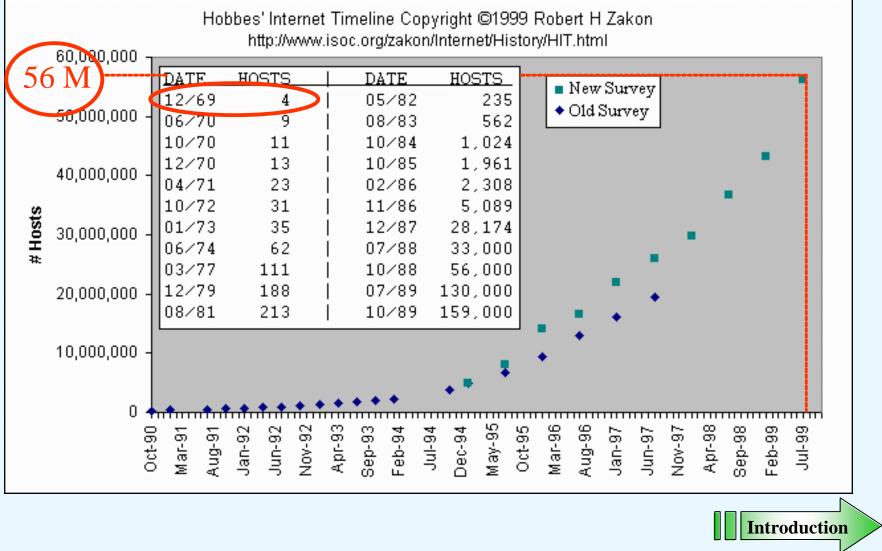
Multimedia Resource Requirements

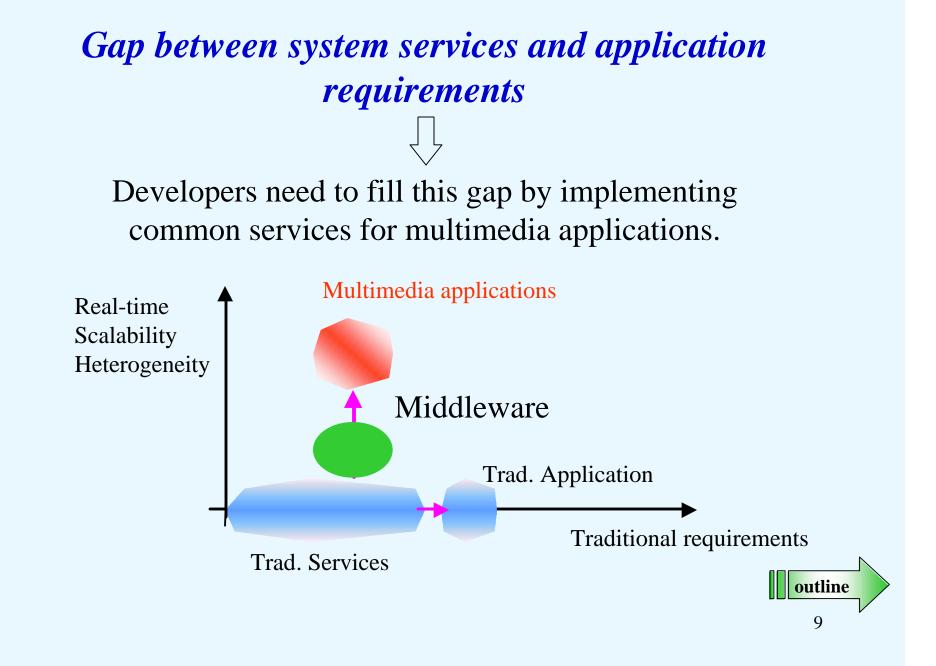


High processing + high bandwidth + Others



Internet Growth





Objective

"Our main objective is to investigate and propose heterogeneous, scalable, reliable, flexible, and reusable solutions and enhancements to common needs in developing multimedia collaborative applications."

Needs we addressed:

- * Extension of network services
- * Media synchronization
- * Floor control
- * Data sharing



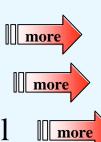
Extension of Network Services

• New services

- * Asynchronous data reception
- * Quality of service monitoring
- * Transmission traffic rate control
- New convenient facilities
 - * Unified Multicast/Unicast API
 - * Efficient buffer management for Application Data Unit II more



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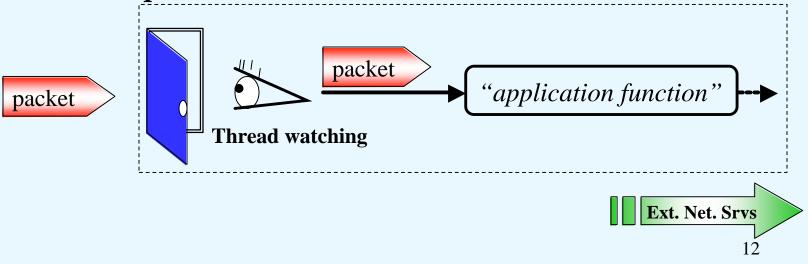
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Asynchronous data reception

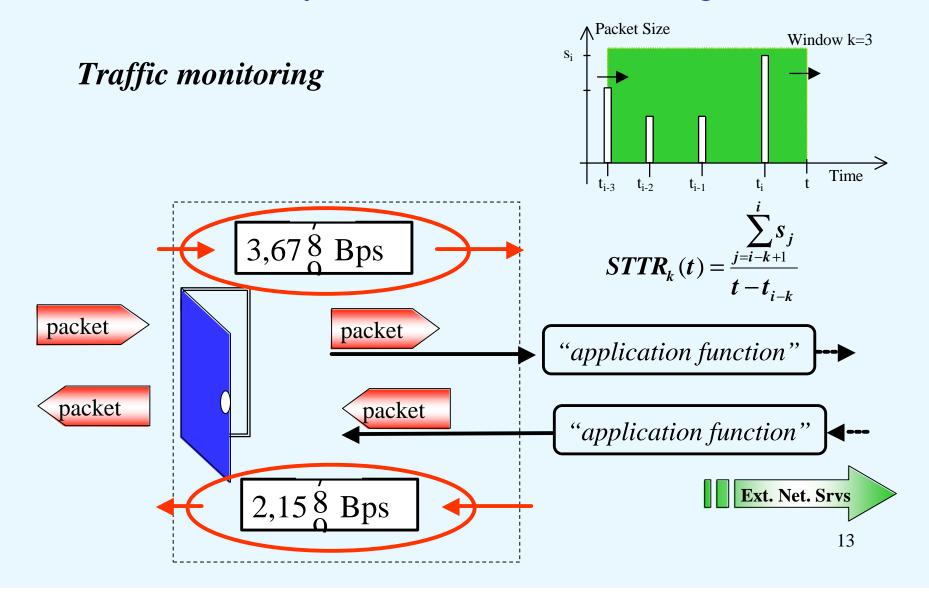
Event-driven model

Defense Outcome	. 🗙
Your Evaluation:	
Pass	Fail

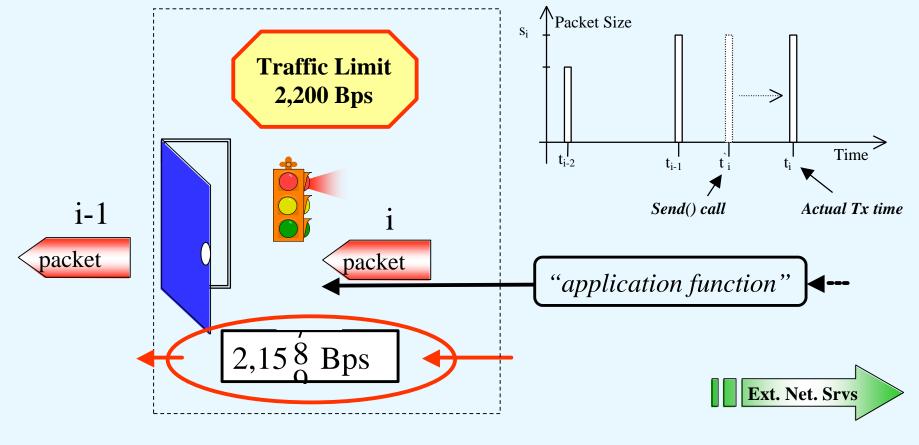
Network packet arrivals



Quality of services monitoring



Transmission traffic rate control



Unified Multicast/Unicast API

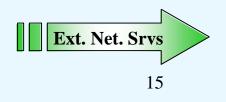
$$\left\{\begin{array}{c} \left\{ Multicast address \\ Unicast address \end{array}\right\}, \quad \left\{ port \right\} \end{array}\right\}$$

• Datagram transmission

* A *send* to a machine or a multicast group does not make a difference.

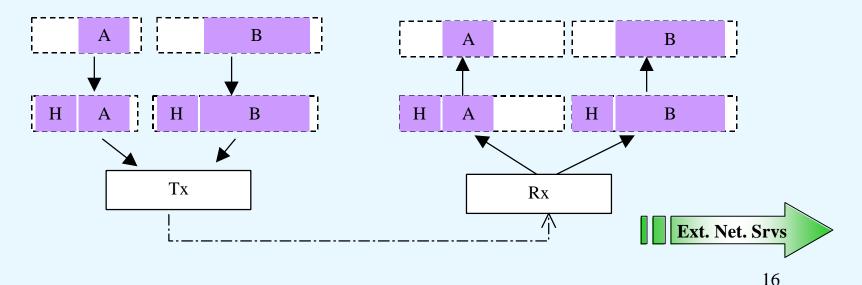
• Datagram reception

- * if the given IP address is a multicast, join group.
- * if address is not multicast, do not bind (I'm client).



Efficient buffer management for Application Data Unit

- * Goal: to prevent payload movements in memory
- * Sender modules create an output buffer that can hold following "headers" and "tails".
- * Receiver module needs to allocate worst case buffer size.



Stream Synchronization

- Problem: processing times and network delays are not deterministic.
- The objective of synchronization is to faithfully reconstruct the temporal relationship between events ("pieces of data").
- Main characteristics of our solution:
 - * It depends on one-way messages only
 - » No need of feedback
 - * It only requires sender's and receivers' clock rates to be constant.
 - » These clocks might be off.
 - » These clocks might even have different rates of change.
 - » No need of globally synchronized clocks
 - * It supports policies to handle late packets and delay adjustments.



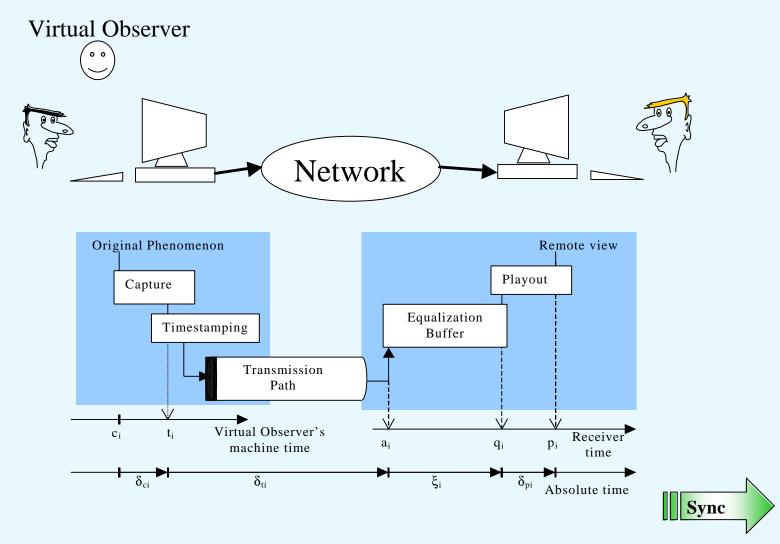
Stream Synchronization (details)

- Time model Imore
- Intra-stream synchronization
- Inter-stream synchronization
- Clock skew estimation and removal



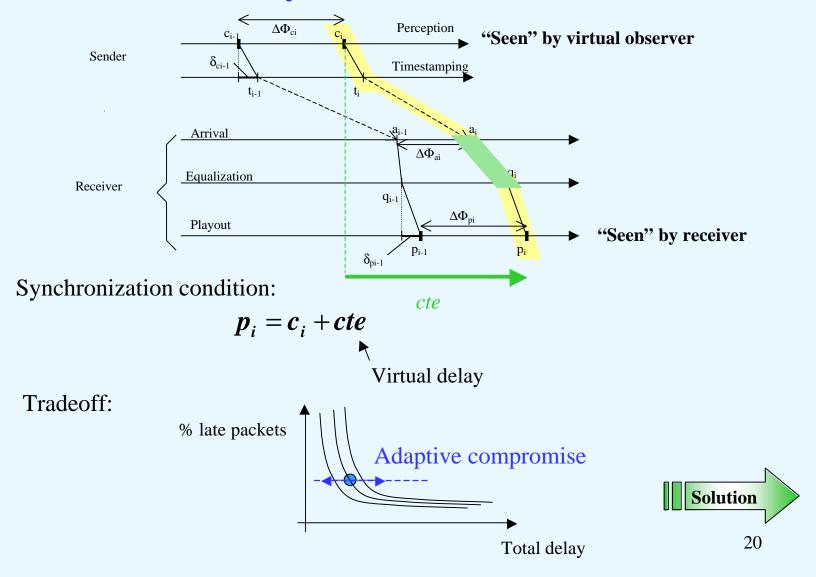


Time Model



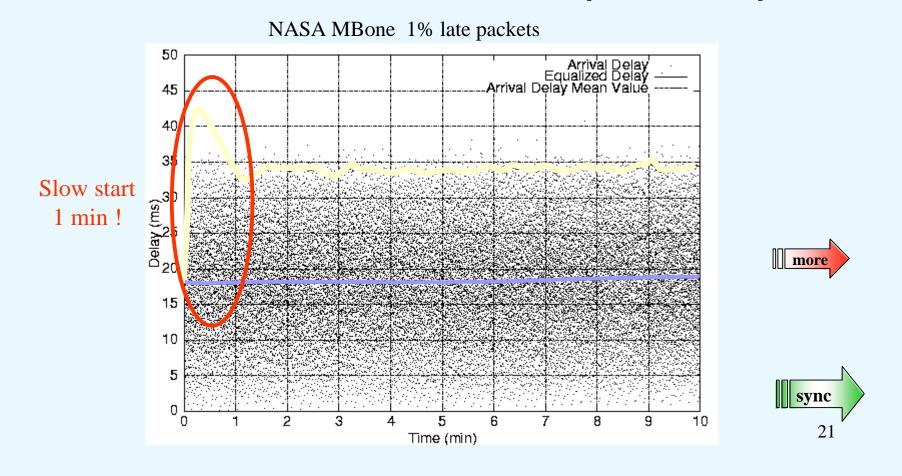
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Intra-stream Synchronization (model)

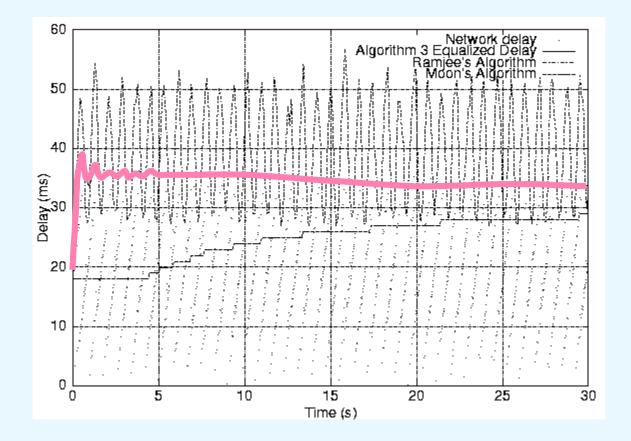


Intra-stream Synchronization (solution)

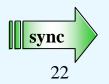
Adjust "virtual delay" to achieve a given % of late packets Estimator for % of late packets: $l_i = a l_{i-1} + (1-a) * \begin{cases} 1 & \text{for late arrival} \\ 0 & \text{otherwise} \end{cases}$



Fast start refinement

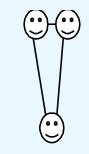


Less than 5 s !



Inter-stream Synchronization

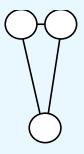
Global synchronization model v/s Differentiated synchronization model



Actual network delay

Global Sync Model

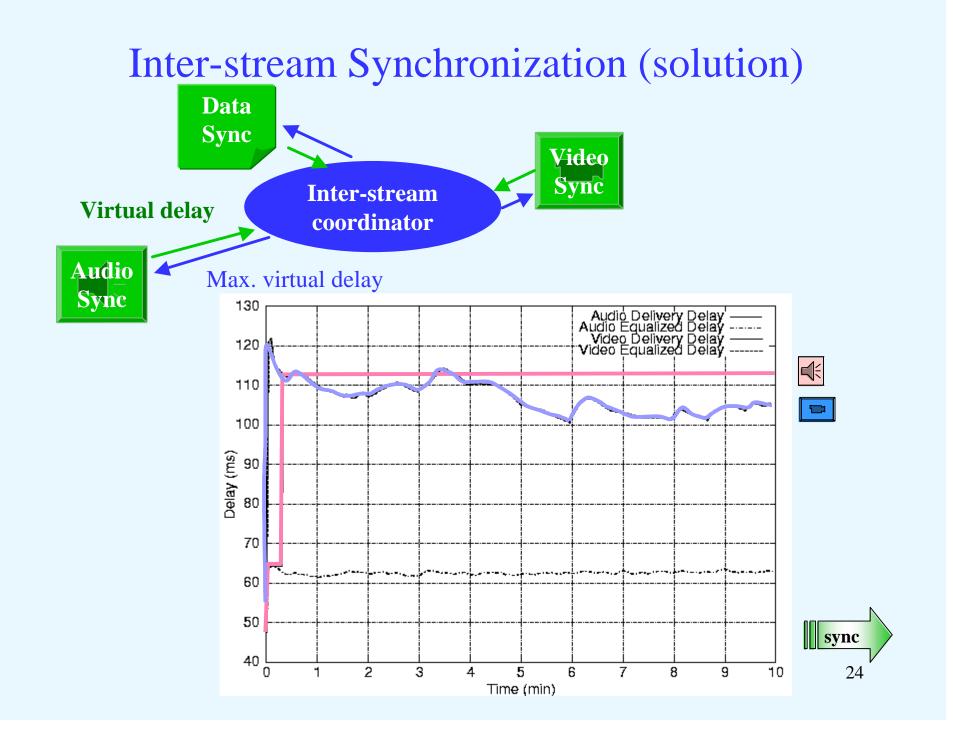
Synchronizes streams coming from anywhere with worst case delay



Differentiated

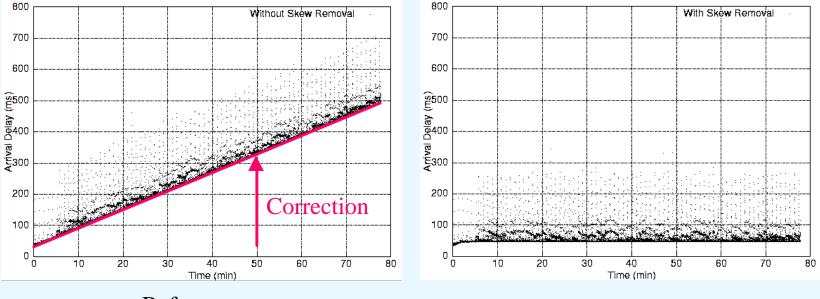
Synchronizes streams coming from one virtual observer





Clock skew estimation and removal

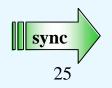
Goal: Remove differences in clock frequencies



Before

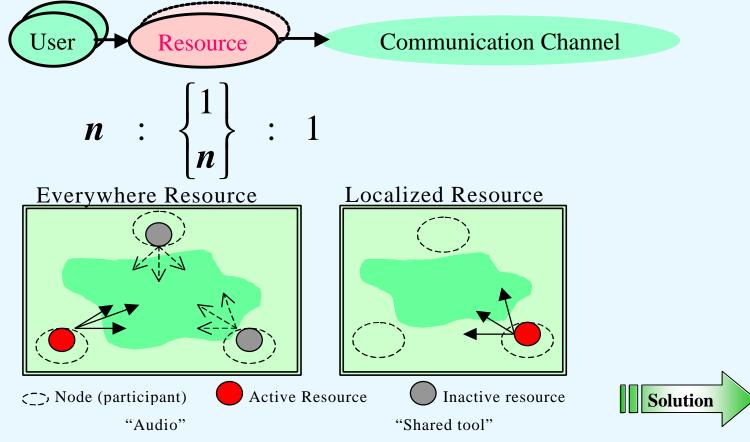
After

The algorithm adjusts a straight line as new packets arrive



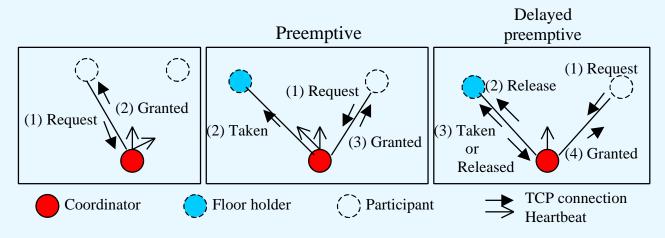
Lightweight Framework for Floor Control

- Problem: How to manage exclusive resources in large-scale multimedia applications?
- We recognize two cases:



Floor Control (Solution)

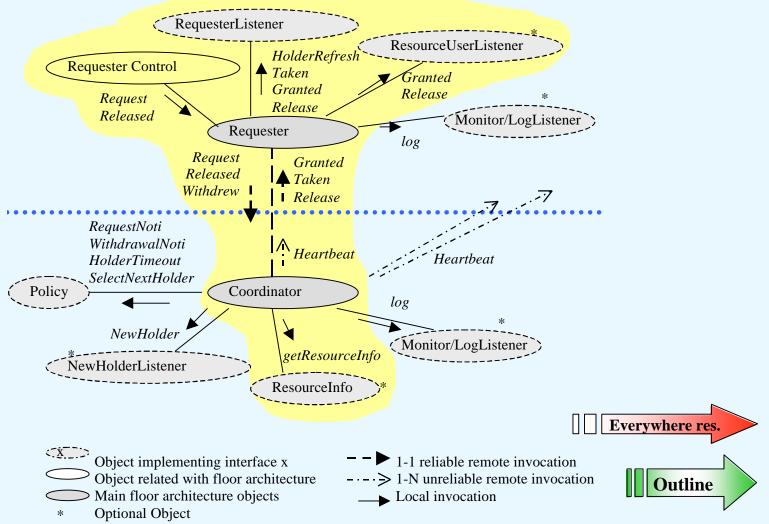
- * We propose two protocols for floor control, one per architecture.
- * Features: lightweight, scalable, robust



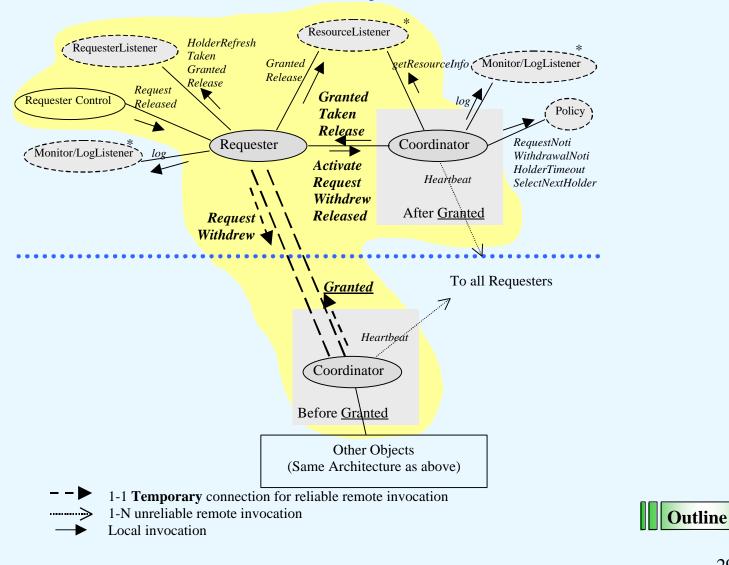
- * The coordinator is stationary for localized resources.
- * The coordinator migrates with floor for everywhere resources



Architecture for localized resources



Architecture for everywhere resources



Protocol for Dynamic Image Transmission

- **Problem**: In addition to audio and video, multimedia sessions needs a component to convey the main idea of discussion.
- Traditional solutions:
 - * Use video (size limitation & high bandwidth)
 - * Shared tools: XTV, co-browsers, VNC,.. (do not scale well)
- Our solution:
 - * Video-like protocol tuned to send dynamic images



Protocol for Dynamic Image Transmission

• Sender:

- * Temporal redundancy removal
 - » Sample image at regular period
 - » Divide image in tiles
 - » Process only changed tiles
- * Spatial redundancy removal
 - » compress and send changed tiles

• Receiver:

- » Receive data unit
- » Decompress tile
- » Update tile in image





Overcoming losses

- Each tile is retransmitted after a random time
- This also accommodates late comers

Performance Study

- * How to select a tile compression technique? (JPEG, GIF, PNG?)
- * Is there a "best" tile size? What does it depend on?
- * How often to sample the image?
- * How can two tiles be compared efficiently?
- * Maximum data transmission rate? What does it depend on?



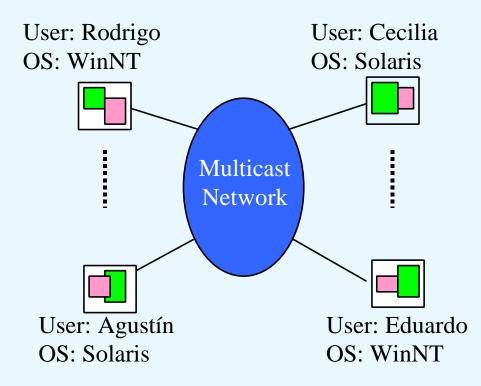
Implementation and Experimental Results

• Implementation:

- * Network support implemented
- * Synchronization: implemented and used with real RTP data in off-line analysis
- * Floor control: partially implemented for localized resources
- * Image protocol implemented
- Putting everything together: Odust
 - * A prototypical sharing tool built on top of the middleware. It uses:
 - * Network support, floor control, dynamic image protocol, other application specific modules.

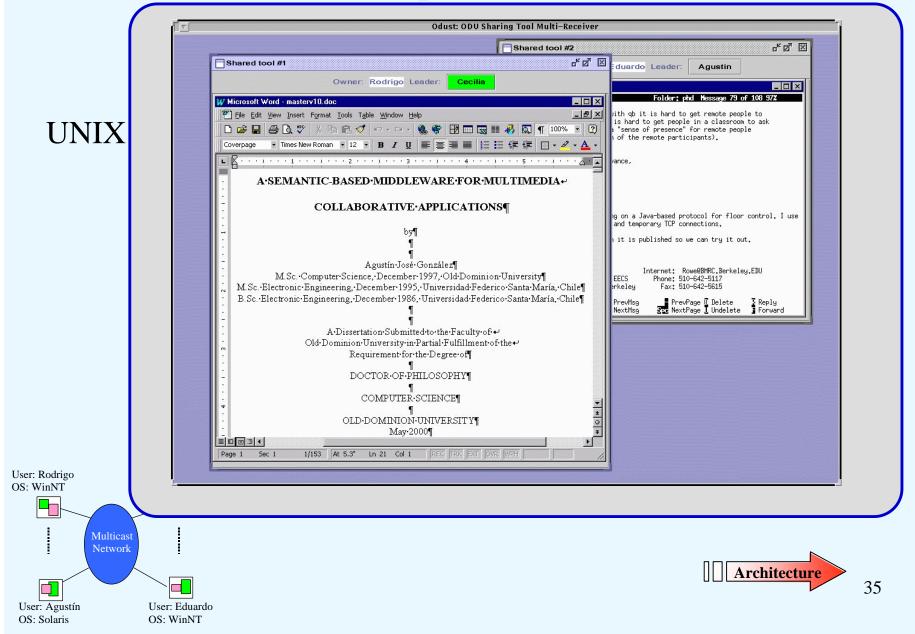


Odust Description

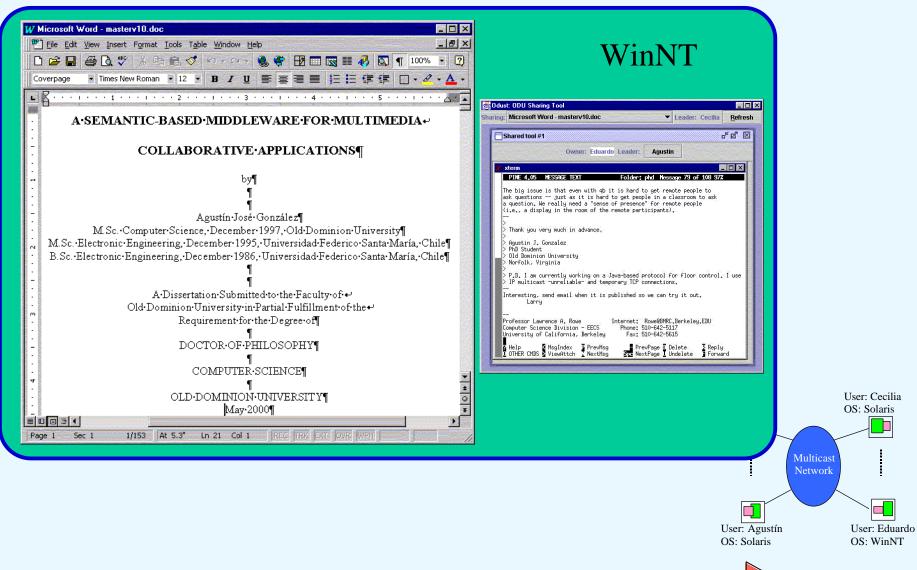




Odust Description: Cecilia's view



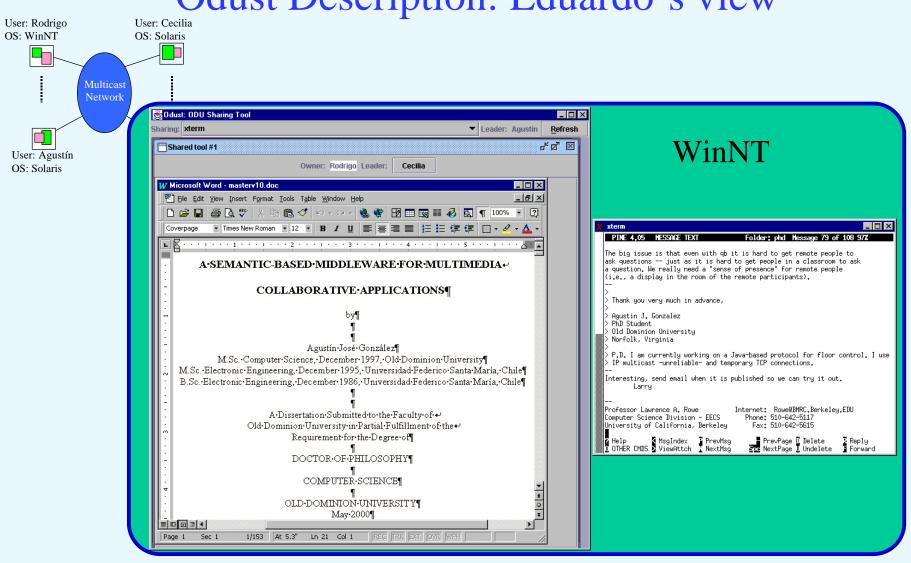
Odust Description: Rodrigo's view





Architecture

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Odust Description: Eduardo's view



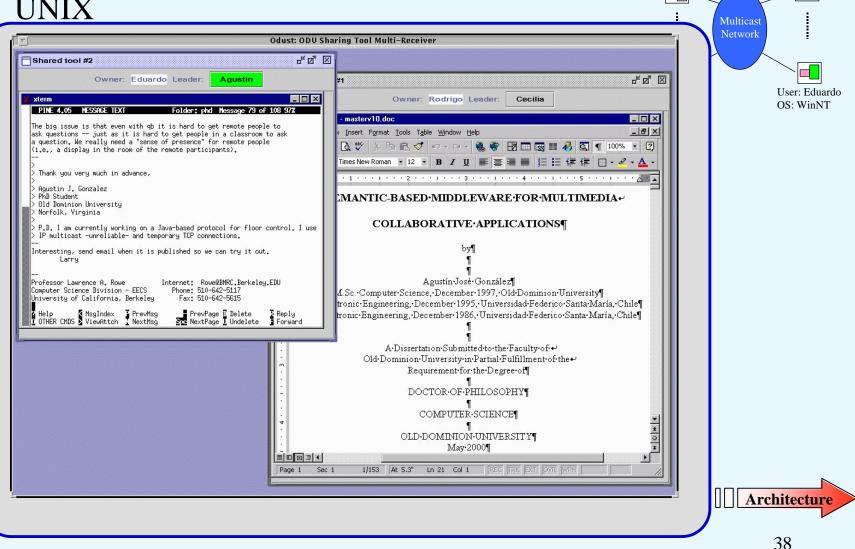
Odust Description: Agustín's view User: Rodrigo

User: Cecilia

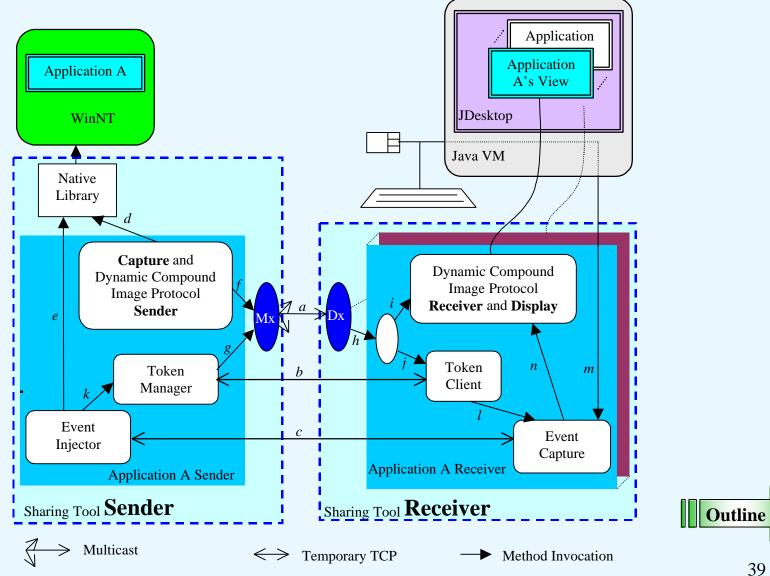
OS: Solaris h

OS: WinNT

UNIX

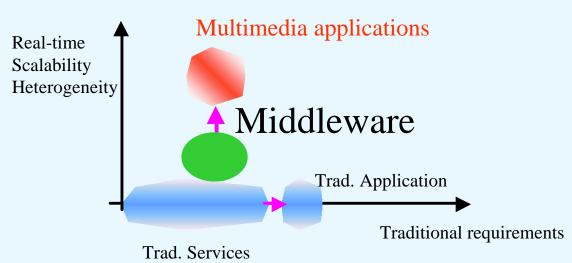


Odust Architecture



Conclusion

• We observed the convenience of a middleware



- It offers:
 - * Multimedia network services
 - * Synchronization
 - * Floor control
 - * Dynamic image transmission

- Future work
 - * Add more components
 - * Continue implementation
 - * Try new ideas (see thesis)

